Minimizing fuel cost in gas transmission networks by dynamic programming and adaptive discretization

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ABSTRACT

In this paper, the problem of computing optimal transportation plans for natural gas by means of compressor stations in pipeline networks is addressed. The non-linear (non-convex) mathematical model considers two types of continuous decision variables: mass flow rate along each arc, and gas pressure level at each node. The problem arises due to the presence of costs incurred when running compressors in order to keep the gas flowing through the system. Hence, the assignment of optimal values to flow and pressure variables such that the total fuel cost is minimized turns out to be essential to the gas industry.

The first contribution from the paper is a solution method based on dynamic programming applied to a discretized version of the problem. By utilizing the concept of a tree decomposition, our approach can handle transmission networks of arbitrary structure, which makes it distinguished from previously suggested methods. The second contribution is a discretization scheme that keeps the computational effort low, even in instances where the running time is sensitive to the size of the mesh. Several computational experiments demonstrate that our methods are superior to a commercially available local optimizer.

1. Introduction

Natural gas has become one of the most important energy resources worldwide. Consequently, the volumes of gas flowing from the fields through transmission networks to the market have been increasing steeply during the past decades, and in parallel, a growing interest in reducing costs associated with pipeline gas transportation has been observed.

A gas transmission network is a system consisting of sources, pipelines, compressors and distribution centers. At the sources, a supply of gas received from external fields is refined, and transmitted via pipelines and compressors to the distribution centers. The distribution centers are the end points of the transmission network, and the gas finally received here is input to local distribution networks supporting households and other clients.

The flow capacity of any pipeline increases with the inlet pressure and decreases with the outlet pressure of the pipeline. If no compressors are installed along a flow path, the pressure will be continuously decreasing. Since the pressure at the distribution centers typically is fixed, the flow capacity may therefore eventually become prohibitively small. To increase the pressure, and thereby the flow capacity, compressors are hence installed at the entry points of selected pipelines. Operation of the compressors incurs a cost depending on the flow and their inlet and outlet pressures.

In this paper, the fuel cost minimization problem (FCMP) to transport natural gas in a general class of transmission networks is addressed. The FCMP involves two types of continuous decision variables: mass flow rate through each arc, and gas pressure level at each node. The problem is to determine a transportation plan minimizing the total fuel cost, while meeting a specified demand at the distribution centers.

An extensive literature on the FCMP has been published over the past 30 years. Most of the suggested solution methods are limited to pipelines networks with acyclic structures, and in such instances, the suggested methods have shown a strong potential. In some of the more recent works, methods for cyclic networks have been developed. However, since these optimization approaches require a certain sparse network structure, their applicability is somewhat restricted. The following sections give a more detailed overview of the most relevant methods. A common assumption is that the system is in steady-state, which means that rapid changes in parameter values do not occur.

1.1. Methods based on dynamic programming

By discretizing the range of the pressure variables, FCMP has in several works been formulated as a combinatorial problem that can be approached by dynamic programming (DP). Wong and Larson (1968) published the first work on optimization of pipeline transportation of natural gas by DP. They applied it to a gun-barrel (linear) network, that is a problem instance where the underlying...
network is a path, using a recursive formulation. A disadvantage was that the length and diameter of the pipeline segment were assumed to be constant because of limitations of DP. March and McCall (1972) modified the problem by adding branches to the pipeline segments and letting the length and diameter of the pipeline segments vary. However, since their problem formulation did not allow unbranched network, more complicated network systems could not be handled.

The first attempt to solve instances with tree-shaped networks by DP was done by Zimmer (1975). A similar approach was described by Lall and Percell (1990). They allowed a divergent branch in their systems and included an integer decision variable into the model that represented the number of operating compressors in the stations.

Carter (1998) developed an algorithm referred to as non-sequential DP. The principal idea of the method is to reduce the network by three basic reductions techniques until it consists of a single node. The method can handle a wide range of instances with cyclic networks, but fails if the networks are not sufficiently sparse. Based on this approach, Borraz-Sánchez and Ríos-Mercado (2004, 2009) developed a hybrid meta-heuristic combing tabu search and non-sequential DP. The restriction that the networks must be sparse is however a shortcoming that the hybrid method inherits from the original paper.

1.2. Methods based on gradient techniques

Percell and Ryan (1987) applied a generalized reduced gradient (GRC) method for solving FCMP. In comparison with DP, an advantage of GRC is that the rapid growth in instance size caused by many discretization points is avoided. Also, GRC is applicable to cyclic networks. Nonetheless, only a local optimum can be provided, of which instances of FCMP can have many, and the solution to be output depends on the choice of starting point. Flores-Villarreal and Ríos-Mercado (2003) extended the previous study by means of an extensive computational evaluation of the GRC method.

1.3. Other techniques and related problems

Wu, Ríos-Mercado, Boyd, and Scott (2000) address the non-convex nature of FCMP, and suggest mathematical models that provide strong relaxations, and hence tight lower bounds on the minimum cost. Based on this model and the PhD thesis of Wu (1998), they demonstrated the existence of a unique solution to a non-linear algebraic equations system over a set of flow variables. This theoretical result lead to a technique for reducing the size of the original network without altering its mathematical structure.


1.4. Contributions from the current work

Several works have demonstrated that, at least in acyclic and sparse cyclic instances of FCMP, it is a promising approach to discretize the pressure variables, and apply DP to the resulting combinatorial problem. The purpose of this research is twofold: First, we demonstrate how such approaches can be applied to networks of arbitrary structure. Second, in order to keep the running time down in dense and cyclic instances, we propose a new scheme for discretizing the pressure variables. This scheme is adaptive in the sense that it avoids fine discretization of variables in area unlikely to contain good solutions, and intensifies discretization in more promising regions.

The remainder of the paper is organized as follows: In the next section, we define the problem in mathematical terms. In Section 3, we present a contemporary solution method, and point out a simple instance where it fails. In Section 4, we show how the weakness of the method discussed in Section 3 can be overcome by our alternative method. Our adaptive discretization method is given in Section 5. Results from computational experiments are reported in Section 6, and concluding remarks are given in Section 7.

2. Problem definition

Let $G = (V,A)$ be a directed graph representing a gas transmission network, where $V$ and $A$ are the node and arc sets, respectively. Let $V^*_v$ and $V^*_v$ denote the sets of out- and in-neighbors, respectively, of node $v \in V$. Let $V_v \subseteq V$ be the set of supply nodes representing the sources, $V_d \subseteq V$ the set of demand nodes representing the distribution centers, and let $A = A_v \cup A_p$ be partitioned into a set of compressor arcs $A_v$ and a set of pipeline arcs $A_p$. That is, if $(u, v) \in A$, then $u, v \in V$ are the network nodes representing the input and the output units, respectively, of some compressor $(u, v)$. An analogous interpretation is made for pipeline arcs $(u, v) \in A_p$.

Two types of decision variables are defined: Let $x_{uv}$ denote the mass flow rate at arc $(u, v) \in A$, and let $p_v$ denote the gas pressure at node $v \in V$. For each $v \in V$, we define the parameters net mass flow rate $b_v$ and pressure bounds $p^l_v$ and $p^u_v$ (lower and upper, respectively). By convention, $b_v > 0$ if $v \in V_v$, $b_v < 0$ if $v \in V_d$, and $b_v = 0$ otherwise. By the assumption that flow is conserved at the nodes, the decision variables are subject to the constraints

$$\sum_{v \in V_u} x_{uv} - \sum_{v \in V_v} x_{uv} = b_v$$

for all $v \in V$. Constraints linking the pressure and flow variables are given for the arc sets $A_v$ and $A_p$, and these are discussed next.

2.1. Compressor arc constraints

The variables that are manipulated in a compressor $(u, v) \in A_v$ in order to have the desired values of $x_{uv}$, $p_u$, and $p_v$ are according to Wu et al. (2000) compressor speed $Su_v$, volumetric inlet flow rate $Q_{in}$, adiabatic head $H_{ad}$ and adiabatic efficiency $\eta_{ad}$ of the compressor. These can briefly be explained as follows (more details can be found in the cited work):

- The variable $Su_v$ is the speed at which each molecule flows through compressor $(u, v)$, and should not be confused with the flow $x_{uv}$ itself.
- While $x_{uv}$ is the mass flow per time unit, the volumetric flow $Q_{in}$ is simply $x_{uv}$ divided by the gas density at the inlet point of the compressor. Due to pressure variations, the density is not constant throughout the network.
- The adiabatic head $H_{ad}$ says how much energy is required to compress one mass unit of gas from one pressure level to another without altering the gas temperature.
- The adiabatic efficiency $\eta_{ad}$ is the ratio between the energy effective in compressing the gas and the total energy spent.

As explained more detailed by, e.g. Wu et al. (2000), the above magnitudes relate to $(x_{uv}, p_u, p_v)$ according to

$$H_{ad} = x^\alpha \left( \frac{p^u_v}{p_u} \right)^{\frac{\kappa}{\kappa - 1}} \quad \forall (u, v) \in A_v$$

(1)

$$Q_{in} = x^\alpha \frac{p^u_v}{p_u} \quad \forall (u, v) \in A_v$$

(2)

$$H_{ad} = \phi^{\alpha} \left( \frac{Q_{in}}{Su_v} \right) \quad \forall (u, v) \in A_v$$

(3)

$$\eta_{ad} = \phi^{\alpha} \left( \frac{Q_{in}}{Su_v} \right) \quad \forall (u, v) \in A_v$$

(4)
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