



# Scheduling a deteriorating maintenance activity and due-window assignment



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## ABSTRACT

Several papers published during the last decade dealt with scheduling a maintenance activity and considered a new setting, where the maintenance duration is assumed to be *deteriorating*, i.e., it requires more time or effort if it is delayed. We study a deteriorating maintenance in the context of due-window assignment, where a time interval is determined such that jobs completed within this interval are “on-time”, whereas early and tardy jobs are penalized. Thus, our paper extends known models by considering *simultaneously* a deteriorating maintenance and due-window. Two deterioration types are considered: *time-dependent* (where the maintenance time increases as a function of its starting time), and *position-dependent* (where it is a function of its position in the sequence). The classical assumption of position-independent processing times was considered first, and then the model is extended to *general position-dependent* processing times. We prove several properties of the optimal timing of the due-window and of the maintenance. Consequently, we show that all the problems studied here are solved in  $O(n^4)$ , where  $n$  is the number of jobs.

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## 1. Introduction

In a recent paper, Yang [21] studied single machine scheduling problems with the option to assign maintenance activities during the production process. The underlying assumption in these models is that, as in many real-life settings, the maintenance activities are *deteriorating*, i.e. they require more time or effort if they are delayed. This new line of research considering deteriorating maintenance activities was introduced first by Kubzin and Strusevich [7]. Following their idea, several other researchers considered various similar settings: Mosheiov and Sidney [14] studied several objective functions such as maximum lateness and number of tardy jobs, Yang [20] studied single machine scheduling problems with both start-time dependent learning and position dependent aging effects, Yang and Yang [23] focused on minimizing the total completion time on a single-machine scheduling with aging/deteriorating effects, Wang et al. [18] investigated parallel machine settings, Cheng et al. [4] and Yang et al. [24] considered a common due window assignment with linear time-dependent deteriorating jobs, and Cheng et al. [3] and Yang et al. [22] focused on deteriorating maintenance on unrelated parallel machines.

In this paper we study a problem of scheduling a maintenance activity on a single machine, with a *due-window assignment*. The basic due-window assignment problem has been introduced and solved by Liman et al. [10], who considered the following setting:  $n$  jobs need to be processed on a single machine around a common due-window; jobs completed within the due-window are “on-time” jobs, whereas early and tardy jobs are penalized. Four cost components were included in the objective function assumed by Liman et al. [10]: earliness, tardiness, due-window starting time and due-window size. They introduced an  $O(n \log n)$  solution, consisting of the job sequence and the determination of the due-window size and location.

Mosheiov and Sarig [12] extended the problem to a setting allowing an option of performing a maintenance activity. On the one hand, this activity requires a fixed (given) time interval during which the machine is turned off, and the following jobs are delayed. On the other hand, the maintenance activity is assumed to be a *rate modifying* activity, i.e. the machine efficiency is improved as a result of the maintenance, and the processing times of the following jobs are shortened. Lee and Leon [8] introduced the first model with a rate-modifying activity. They were followed by Lee and Lin [9], Kubzin and Strusevich [7], Gordon and Tarasevich [5], Zhao et al. [26], Lodree and Geiger [11], Jin and Cheng [6], and Zhao and Tang [25], Wang and Wang [19], among others. The problems solved in these papers are clearly harder, as they require in addition to the

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traditional job scheduling and due-window assignment decisions, a decision regarding the optimal timing of the maintenance activity.

In this paper we extend the above models by combining a due-window assignment and the option of scheduling a *deteriorating* maintenance activity. In this setting the maintenance time increases if it is delayed. As written in Mosheiov and Sidney [14], “the longer the maintenance activity is delayed, the worse the system conditions become, so that the maintenance requires more effort and time”. We assume two deterioration types: *time-dependent* deterioration and *position-dependent* deterioration. First, as in Kubzin and Strusevich [7], we assume *linear* deterioration, i.e. the maintenance time increases linearly with its starting time. Then, we assume that the deterioration time increases (in the most general form) as a function of its position. Consider the following typical example of a car maintenance: the driver is required to go through a maintenance procedure either after a certain amount of time (reflecting time dependent deterioration), or a specific distance traveled (position-dependent deterioration). We show that in both cases the extension does not increase the complexity of the problem (compared to the associated due-date assignment model). We further extend these models to allow *general position-dependent* job processing times. These extensions do not increase the total complexity as well, i.e. the problems are solved in  $O(n^4)$  time, where  $n$  is the number of jobs. Several recent papers appear to be particularly relevant. First, we note that Cheng et al. [4] and Yang et al. [24] studied due-window assignment with a deteriorating maintenance as well. However, in their models: (i) the job processing times are deteriorating with *job-independent* deterioration parameters, and (ii) after performing the maintenance, the machine reverts to its initial condition. Their problems are shown to be solved in  $O(n^2 \log n)$  time. The maintenance considered in our model is, as mentioned, a *rate modifying* activity, and the effect on the job processing times is *job-dependent*. Rustogi and Strusevich [15] studied a more general model, in which the scheduler may perform more than a single maintenance activity. Unlike previous papers dealing with scheduling multi-maintenance activities, Rustogi and Strusevich [15] considered a very realistic assumption that the maintenance does not necessarily fully restore the machine to its original state. We note that this model (i) does not consider a due-window (the objective function is minimum makespan, whereas our objective is minimum earliness, tardiness and due-window costs), and (ii) assumes only *deteriorating* job processing times, where we allow *general position dependent* processing times (and no monotonicity is enforced). Rustogi and Strusevich [16] studied a very general scheduling model, where the contribution of a job to the objective function is given by a product of its processing time and a certain positional weight. It is well known that many scheduling problems fall into this category and, depending on the input matrix, can either be reduced to a linear assignment problem, or even be solved by simple matching. The input matrix considered in their paper is assumed to be determined by a product of two arrays (i.e.,  $C_{ij} = \alpha_i \beta_j$  for the given arrays  $(\alpha_1, \alpha_2, \dots, \alpha_n)$  and  $(\beta_1, \beta_2, \dots, \beta_n)$ ). It should be noted that the input matrix assumed in our paper (for the case of position-dependent processing times) is *not restricted* in any way. Finally, Rustogi and Strusevich [17] studied another extension of a model dealing with a single machine and several (potential) maintenance activities, in which the actual processing times of the jobs are subject to a combination of positional and time-dependent effects. However, unlike the model studied here, these effects in Rustogi and Strusevich [17] are *job-independent*. Thus, in our knowledge, the problem combining (i) a due-window assignment, (ii) scheduling a deteriorating maintenance, and (iii) general position-dependent job processing times, is studied here for the first time.

In Section 2 we provide the notation and the formulation of the problem. Section 3 summarizes a list of properties of an optimal schedule. Sections 4 and 5 present the solutions for the cases of

time-dependent deterioration and of position-dependent deterioration, respectively.

## 2. Formulation

We study an  $n$ -job, single-machine scheduling problem. An option of scheduling a maintenance activity (rate modifying) may be considered by the scheduler. First, as in Lee and Leon [8], we assume that the length of the maintenance is a function of its starting time. If it starts at time  $t$ , then its length is  $T_{MA}(t) = t_0 + \omega t$ , where  $t_0 > 0$  is its basic maintenance time (if performed at time zero), and  $\omega > 0$  is a non-negative deterioration factor. The case of a linear decreasing maintenance time (learning) is also considered:  $T_{MA}(t) = t_0 - \omega t$ . Then, we assume that the maintenance is a function of its position: if it is scheduled before the job in position  $r$ , it requires  $T_{MA}(r) = \tau_r$  time, where  $\tau_r$  is a position-dependent constant, and  $\tau_r \leq \tau_{r+1}$ ,  $r = 1, \dots, n-1$ . During the maintenance no production is performed. In the simplest model studied here, the processing times of the jobs scheduled prior to the maintenance are assumed to be position-independent, and similarly, the processing times of the jobs scheduled after the maintenance are assumed to be position-independent. Specifically, the processing time of job  $j$  is  $P_j$  if the job is processed prior to the maintenance activity, and  $\theta_j P_j$  ( $0 < \theta_j \leq 1$ ) if it is scheduled after it,  $j = 1, \dots, n$ .  $\theta_j$  is the *modifying rate* of job  $j$ . We also use the notation  $P_r$  and  $\theta_r$  to denote the processing time of the job in the  $r$ -th position, and the modifying rate of the job in the  $r$ -th position,  $r = 1, \dots, n$ , respectively. This basic model is extended to allow position-dependent processing times. In this case  $P_{jr}$  denotes the processing time of job  $j$  if assigned to position  $r$ ,  $j, r = 1, \dots, n$ .

For a given job sequence, the completion time of the job in position  $r$  is denoted by  $C_r$ ,  $r = 1, \dots, n$ . All the jobs share a common due-window:  $[d_1, d_2]$ , such that  $d_1 \leq d_2$ . The earliness of the job in the  $r$ -th position is given by  $E_r = \max\{0, d_1 - C_r\}$  and the tardiness of the job in the  $r$ -th position is given by  $T_r = \max\{0, C_r - d_2\}$ ,  $r = 1, \dots, n$ . The objective function consists of four cost components:  $\alpha$  is the earliness unit cost,  $\beta$  is the tardiness unit cost,  $\gamma$  denotes the unit cost of (delaying) the due-window starting time, and  $\delta$  is the unit cost of (increasing) the due-window size. Let  $D = d_2 - d_1$  denote the due-window size. Then the objective function is (see [10]):

$$Z = \sum_{r=1}^n (\alpha E_r + \beta T_r + \gamma d_1 + \delta D).$$

## 3. Basic properties: due-window and maintenance location

We quote a number of properties of an optimal schedule proved in Mosheiov and Sarig [12] for the case of a *fixed* maintenance time. These properties are easily shown to be valid for our more general setting of deteriorating maintenance.

**Property 1.** An optimal schedule starts at time zero and contains no idle time between consecutive jobs.

**Property 2.** An optimal schedule exists in which the maintenance activity is performed prior to the starting time of one of the jobs, or not performed at all.

Define:

Condition (i):  $\gamma > \delta$ .

Condition (ii):  $\beta < \min\{\gamma, \delta\}$ .

**Property 3.1.** If condition (i) holds, an optimal schedule exists in which the due-window starts at time zero.

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