



A genetic algorithm approach for a single hoist scheduling problem with time windows constraints



Adnen El Amraoui^{a,b,*}, Marie-Ange Manier^a, Abdellah El Moudni^a, Mohamed Benrejeb^b

^a IRTES-SET, Department of Engineering and Management of Process, Université de Technologie de Belfort-Montbéliard, UTBM, 90010 Belfort, France

^b Department of Electrical Engineering, Ecole Nationale d'Ingénieurs de Tunis, ENIT, BP 37, le Belvédère, 1002 Tunis, Tunisia

ARTICLE INFO

Article history:

Received 4 June 2012

Received in revised form

17 December 2012

Accepted 16 February 2013

Available online 24 April 2013

Keywords:

Cyclic hoist scheduling problem
Electroplating industry
Genetic Algorithm (GA) approach
Multi-job
Time window constraints
Transportation resource

ABSTRACT

This study considers a cyclic scheduling of hoist movements in electroplating industry. Several jobs have to flow through a production line according to an ordered bath sequence. They firstly enter the line at a loading buffer. Then they will be soaked sequentially in a series of tanks containing specific chemical baths. Finally, they will leave the line at the unloading buffer. The processing time duration of each job in each tank is not constant but confined within a time window bounded by a minimum and a maximum duration. If a job spends less than the minimum duration or more than the maximum duration it is considered defective. Moreover, not only the job operations in the soaking tanks have to be scheduled, but also the transportation of the jobs between tanks has to be considered. The problem now is to find an optimum or near optimum feasible cyclic scheduling such that the hard resource and time-window constraints are respected and the cycle time duration is minimized. A mathematical formulation is proposed for the multi-jobs cyclic hoist scheduling problem with a single transportation resource, and a Genetic Algorithm (GA) approach is presented to solve it. The performance of the proposed algorithm is evaluated with the objective value obtained with a linear programming model, on several literature instances. Computational experiments show the good performance of our GA in terms of solution quality, convergence and computation time.

© 2013 Elsevier Ltd. All rights reserved.

1. Introduction

In recent years scheduling problem with transportation resources has gained much attention in the literature. Despite the differences between the several proposed models, they all deal with the coordination of material handling devices and classical problem scheduling decisions. In this section we review the most important research concerning the Cyclic Hoist Scheduling Problem (CHSP). This problem is a branch stemming of the Hoist Scheduling Problem (HSP). Since the first model given by Phillips and Unger (1976), several researchers have been interested in the Hoist Scheduling Problem (Manier and Bloch, 2003) and a large number of mathematical models (Spacek et al., 1999; Chen et al., 1998; Zhou and Ling, 2003; Che and Chu, 2007) have been developed. Nevertheless, the cyclic problem and more precisely, the single-job cyclic schedule, was the most studied case in literature (Armstrong et al., 1996; Manier et al., 2000) due to the strong complexity of the problem, even with a simple line configuration (single-job, single-hoist and single-capacity

tanks). However, the problem is so complex that the most practical path consists of addressing the simplest cases, and then studying possible extensions. One of these extensions that we propose to study in this paper is the multi-parts jobs. For a given cyclic schedule, we have to minimize the cycle time duration for n different part-jobs. Each cycle is specified by a parameter n , called the cycle degree. In an n -degree cyclic schedule, exactly n parts enter and leave the production line every one cycle (Lei, 1994). The most known studies dealing with this extension are presented as follows.

Ptuskin studied a multi-parts problem (single job problem, line with 50 tanks, n parts jobs and n operations) denoted (CHSP|50/nc| n parts/ n ps, n op| C_{max}) according to the notation, elaborated by Manier and Bloch (2003) and dedicated to Hoist Scheduling Problems. In this problem, parts are processed according to the same sequence, with various processing times, where the sequence periodically alternates jobs of different parts and is known in advance. Moreover, the date of each part has to be computed. This problem is decomposed in several mono-product sub-problems and the solution corresponds to a common period (Ptuskin, 1995). Nevertheless, the obtained cyclic schedule is feasible but it is not guaranty to be optimal.

Varnier and Jeunehomme used a branch and bound method for a similar multi-part problem (CHSP|mt//load-unload| n ps, n op| T_{min} problem). A difference with Ptuskin (1995) is that the entry

* Corresponding author at: Department of Engineering and Management of Process, Université de Technologie de Belfort-Montbéliard, UTBM, 90010 Belfort, France. Tel.: +33 384583796.

E-mail addresses: elamraoui.adnen@gmail.com, Adnen.ElAmraoui@enit.rnu.tn, adnen.elamraoui@utbm.fr (A. El Amraoui).

sequence of jobs (configuration) is not known in advance. Each level of the search tree consists in adding a possible state at the beginning of the cycle, which is called configuration. The possible schedules from this configuration constitute the branches of the search tree and then a linear program is achieved to check if the problem has a solution. If not, backtrack is allowed to consider the other schedules. However, Varnier and Jeunehomme mainly study the transition part of the schedule (Varnier and Jeunehomme, 2000).

A branch and bound method has also been used by Mateo and Companys to study the CHSP $|mt|/diss|/2,2mt+2|T_{min}$ for a two different part jobs problem. The branch and bound procedure builds the sequence of movements progressively. Each level of the search tree consists of adding one tank and thus, the stages to be done on it (stage is a new indexation used by authors to define the soaking operations). A linear program is then solved at each node to check the consistency of the constraint system. Nevertheless, this method cannot be easily extended to the general n -parts jobs (Mateo and Companys, 2006).

Lei and Liu present first, a formal analysis of the hoist scheduling problem with two non-identical parts (CHSP $|1,20,ct|/|42|T_{min}$). Then, the results are used to develop a new branch-and-bound procedure in order to find the optimal schedules. This procedure uses a depth search strategy. In each step, it chooses a node at the greatest depth to branch first. Whenever a tie is found, it chooses the one with the minimum lower bound to branch first. They state that each node at level k of their tree has two partial cycles. Then the minimum lower bound given by those two cycles is used as the lower bound of this node. For each newly generated node, authors use the result of their analysis to check its feasibility. If the offspring nodes are either infeasible or do not improve the best obtained lower bound, the node is eliminated and backtrack is allowed to its upper level. This procedure provides good computational results. Authors tested their procedure on mono-product benchmarks. They also randomly generated some test problems without really detailing them. But the efficiency and robustness of their procedure was not evaluated considering variations of hoist speed and width of time windows. Moreover, their model was dedicated to basic line configurations' (single capacity resources ...) and for 2-degree cyclic schedule (Lei and Liu, 2001).

Many other studies deal with scheduling of the other production lines or even with the CHSP with relaxed constraints: mainly bounded processing times and no-wait constraints; which are problems near the CHSP (Fleury et al., 2001; Ge and Yih, 1995; Che et al., 2009; Che and Chu, 2005; Kats et al., 2006; Liu et al., 2002; Song et al., 1995).

A synthesized comparison between these works (dealing with the Multi-Parts CHSP) is presented in Table 1. Therefore, all of these works carried out on this issue, to date, have several limitations. Thus and in order to improve the electroplating line production, a general n -cyclic model could be developed by incorporating additional constraints to previous works.

Getting motivated by such class of scheduling problem and previous researches using evolutionary approach, we propose a genetic algorithm approach in order to optimality or

near-optimality solve an n -degree CHSP problem for n different jobs. In such kind of industrial problems, electroplating companies are looking for optimizing the simultaneous manufacturing of several kinds of products through a single line. This paper can be considered as the first paper tackling the n -degree cyclic hoist scheduling problem with heterogeneous part-jobs and time-window constraints. And it is organized as follows. In Section 2, we formally describe the problem and formulate it as a Mixed Integer Linear Programming model. In Section 3, we firstly propose a conjunctive graph model of the problem and then a polynomial computational procedure to compute the cycle time for a fixed cyclic schedule. In Section 4, we develop our Genetic Algorithm approach. At last, in Section 5 and Section 6, we make some concluding remarks based on the computational results and suggest directions for future research.

2. Problem statements

We consider the single-hoist cyclic scheduling problem in an automated system (CHSP $|mt|/diss|/n,nop|T_{min}$) which may be described as follows. We are given mc tanks M_1, \dots, M_{mc} and n jobs J_1, \dots, J_n . In general, loading and unloading buffers (or tanks) are assimilated to fictitious tanks numbered respectively 0 and $mc+1$. Processing stations are containing chemical baths and arranged in a row (see Fig. 1). Each job J_j consists of n_j operation treatments O_{ij} $\{i=1, \dots, n_j\}$. During the soaking process, no preemption is allowed. It means that the soaking operation O_{ij} has to be processed in one time on a dedicated tank $\alpha_{ij} \in \{M_1, \dots, M_{mc}\}$ for $p_{ij} \geq 0$ time units (t.u.). This soaking duration, is confined within a minimum ($0 \leq a_{ij} \leq p_{ij}$) and a maximum ($p_{ij} \leq b_{ij} \leq +\infty$) time durations as specified in the processing specifications. One of the specificities of the CHSP is that consecutive soaking operations of the same job are not allowed to be assigned to the same tank. Moreover, we assume that tanks cannot receive more than one job on same time, in this case we speak about what we call a single capacity tank ($ct=1$). Precedence constraints are given by the job processing sequences. Where, $O_{ij} \rightarrow O_{i+1j}$ means that the i th soaking operation of job J_j has to be finished completely before starting $i+1$ th soaking operation of the same job. This problem description can be valuable for a general shop-scheduling problem. Nevertheless, the CHSP presents additional specified constraints class, called transport times constraints. They occur when a job has to be transported from one tank to another, i.e. if a job J_j has to be transported from tank M_u to M_v , the hoist (or automated guided vehicles) insuring this moving operation ($m_{j,u,v}$) has to spend $d_{u,v}$ t.u. This duration can be, sometimes, job dependent ($d_{j,u,v}$). We suppose that the hoist has a single job capacity, where it cannot handle more than one job at the same time. Moreover, as different from the classical shop problems, it is not allowed that jobs have to wait before to finish all there soaking operations. This means that no buffer space exists between tanks and the hoist is not allowed to wait while transporting a job.

In this paper, we consider that all jobs are processed on the same order soaking tanks, but with different time windows-soaking

Table 1
Comparison between different studies dealing with the Multi-Parts CHSP.

Problem	TW	CD	PS	Objectif	Method	Optimality	Problem specificities
Ptuskin (Ptuskin, 1995)	Identical	n	Single	C_{max}	Fuzzy Logic	No	Feasible solution
Varnier (Varnier and Jeunehomme, 2000)	Different	3	Severel	T	Branch & Bound	Yes	Transitional schedule
Mateo (Mateo and Companys, 2006)	Different	2	Single	T	Branch & Bound	Yes	Basic problem
Lei (Lei and Liu, 2001)	Different	2	Single	T	Branch & Bound	Unknown	Formal analysis

TW: Time Window, CD: Cyclic Degree, PS: Processing Sequence

متن کامل مقاله

دریافت فوری ←

ISIArticles

مرجع مقالات تخصصی ایران

- ✓ امکان دانلود نسخه تمام متن مقالات انگلیسی
- ✓ امکان دانلود نسخه ترجمه شده مقالات
- ✓ پذیرش سفارش ترجمه تخصصی
- ✓ امکان جستجو در آرشیو جامعی از صدها موضوع و هزاران مقاله
- ✓ امکان دانلود رایگان ۲ صفحه اول هر مقاله
- ✓ امکان پرداخت اینترنتی با کلیه کارت های عضو شتاب
- ✓ دانلود فوری مقاله پس از پرداخت آنلاین
- ✓ پشتیبانی کامل خرید با بهره مندی از سیستم هوشمند رهگیری سفارشات